

Figure 16: Topology for two sources and two consumers, problem 73.

(a) Consider a network that supplies a good (e.g. water, electricity or cookies) to consumers. The network has two suppliers and two consumers. Source *i* produces quantity s_i and supplies a fraction f_i to consumer *i*, and a fraction $1-f_i$ to consumer *j*, as in fig. 16. Source *j* acts similarly. Each consumer consumes whatever is supplied. Thus the consumption by consumer *i* is:

$$c_i = f_i s_i + (1 - f_j) s_j \tag{299}$$

where i = 1, 2 and j = 3 - i.

There is fractional-error uncertainty in the source properties:

$$\mathcal{U}(h) = \left\{ f_i, s_i : \left| \frac{f_i - \tilde{f}_i}{\tilde{f}_i} \right| \le h, \ f_i \in [0, 1], \left| \frac{s_i - \tilde{s}_i}{\tilde{s}_i} \right| \le h, \ s_i \ge 0, \ i = 1, 2 \right\}, \quad h \ge 0$$
(300)

We require that each consumer be within δ of a specified value, \overline{c} :

$$|c_i - \overline{c}| \le \delta \tag{301}$$

Derive an explicit expression for the inverse of the robustness function for consumer *i*. Evaluate and compare the robustnesses for $\overline{c} = 1$, $\tilde{f} = 1/2$ and $\tilde{s} = 0.9$ or 1.0. Which option is preferred? Why, and what does this mean?

(b) Modify part 73a as follows. The nominal consumption by each consumer is \tilde{c} but the actual consumption is:

$$c = \tilde{c} + \varepsilon \tag{302}$$

where ε is an exponentially distributed random variable whose pdf is $p(\varepsilon) = \lambda e^{-\lambda \varepsilon}$, $\varepsilon \ge 0$. Derive an explicit algebraic expression for the probability that $c \le \overline{c}$ where \overline{c} is a known positive value greater than \tilde{c} .

(c) Continuing part 73b, consider uncertainty in the exponential coefficient, λ , represented by the info-gap model:

$$\mathcal{U}(h) = \left\{ \lambda : \ \lambda > 0, \ \left| \frac{\lambda - \widetilde{\lambda}}{w} \right| \le h \right\}, \quad h \ge 0$$
(303)

where $\tilde{\lambda}$ and w are known positive values. We require that the probability that $c \leq \overline{c}$ be less than δ where $0 < \delta < 1$. Derive an explicit algebraic expression for the robustness.

- (d) Continuing part 73c, consider two different designs with values λ_i and w_i , for i = 1 and 2. Based on the robustness function, derive an explicit algebraic expression for the values of δ for which you prefer system 1.
- (e) A particular consumer (you, perhaps) is supplied by N sources resulting in consumption equal to:

$$c = \sum_{i=1}^{N} f_i s_i \tag{304}$$

The fractions f_i and source terms s_i are uncertain:

$$\mathcal{U}(h) = \left\{ f, s: f_i \ge 0, \left| \frac{f_i - \tilde{f}_i}{\tilde{f}_i} \right| \le h, \ s_i \ge 0, \left| \frac{s_i - \tilde{s}_i}{\tilde{s}_i} \right| \le h, i = 1, \dots, N \right\}, \quad h \ge 0$$
(305)

We require that the consumption be no less than the critical value \overline{c} :

$$c \ge \overline{c}$$
 (306)

Derive an explicit algebraic expression for the robustness function.

(f) Repeat part 73e with the modification that the fractions, f_i , are positive and known for sure and the source terms are uncertain according to:

$$\mathcal{U}(h) = \left\{ s : (s - \tilde{s})^T W^{-1} (s - \tilde{s}) \le h^2 \right\}, \quad h \ge 0$$
(307)

where \tilde{s} is a known vector and W is a known, positive definite, real, symmetric matrix. The performance requirement is eq.(306). Derive an explicit algebraic expression for the robustness function.

(g) Consider the following modification of the 2-source and 2-consumer network, in which we introduce a mutual commitment. Under ordinary conditions, each source supplies a single consumer at each discrete time step:

$$c_1(t) = s_1, \quad c_2(t) = s_2, \quad t = 0, 1, 2, \dots$$
 (308)

Each consumer has its own private supply, and each consumer requires a positive consumption \overline{c}_i , i = 1, 2.

However, the consumers have mutual commitments. If consumer *i* loses its supply at some time step, then in the next time step consumer *j* is committed to supply consumer *i* with a fraction γ of *i*'s requirement, \overline{c}_i , though *j* cannot supply more than s_j provides.

Suppose that at some time step, call it t = 0, consumer 1 loses its supply, so the consumption in this step is:

$$c_1(0) = 0, \quad c_2(0) = s_2$$
 (309)

In the next time step, consumer 2 must transfer to consumer 1 a part of 2's supply so the consumption is:

$$c_1(1) = \min(s_2, \gamma \overline{c}_1), \quad c_2(1) = (s_2 - \gamma \overline{c}_1)^+$$
 (310)

where $x^+ = x$ if $x \ge 0$ and equals zero otherwise.

The sources are uncertain according to a fractional-error info-gap model:

$$\mathcal{U}(h) = \left\{ s: \ s_i \ge 0, \ \left| \frac{s_i - \widetilde{s}_i}{w_i} \right| \le h, i = 1, 2 \right\}, \quad h \ge 0$$
(311)

where \tilde{s}_i and w_i are known positive constants.

Derive an explicit algebraic expression for the robustness of consumer 2 at step 1.

(h) Using the robustness function from part 73g, consider the following two commitment situations, (γ, w_2) and (γ', w_2') , where:

$$\gamma' < \gamma, \quad w_2' > w_2 \tag{312}$$

The 'prime' configuration entails lower commitment by consumer 2, but greater uncertainty in consumer 2's source. Use the robustness function to discuss the values of consumer 2's required consumption, \overline{c}_2 , for which 2 prefers the 'prime' configuration.